
Watersheding hierarchies

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Hierarchical watersheds

- Sequences of nested partitions such that each partition is a watershed of a filtering of the initial relief/map according to regional attributes, e.g., on geometric, photometric, or learned information
- Hierarchies of segmentation can be equivalently represented by saliency maps

Original image



Hierarchical watershed \mathcal{H}



Saliency map of \mathcal{H}



- Hierarchical watersheds are linked to a broader family of combinatorial optimization problems: minimum spanning trees, random walkers and graph cuts [C09]
- They satisfy interesting mathematical properties (e.g. they preserve a minimum spanning tree of the original graph)
- They are useful in object detection and perform well on natural images [P18]

[C09] C. Couprie, L. Grady, L. Najman, and H. Talbot. *Power watersheds: A new image segmentation framework extending graph cuts, random walker and optimal spanning forest.*

[P18] B. Perret, J Cousty, S. Guimaraes, D. Maia. *Evaluation of hierarchical watersheds.*

Problem

Given a hierarchy \mathcal{H} and a weighted graph (G, w)

- *find a hierarchical watershed \mathcal{H}' of (G, w) which “approximates” \mathcal{H}*

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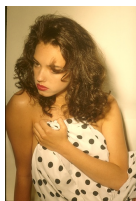
Contributions

Introduction of the watershedding operator

- *which transforms any hierarchy into a hierarchical watershed, and*
- *whose fixed points are the hierarchical watersheds*

Watershedding: applications

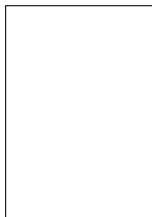
- Refinement of coarse hierarchies
 - e.g., high quality hierarchical segmentation method (COB) [K18]



Original



SED



COB



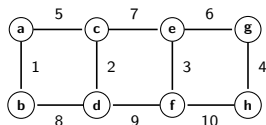
Watershedding

[K18] K.-K. Maninis, J. Pont-Tuset, P. Arbeláez, and L. Van Gool. *Convolutional oriented boundaries: From image segmentation to high-level tasks*. *PAMI*. 2018.

- 1 Hierarchical watersheds
- 2 Watersheding transform
- 3 Conclusion and perspectives

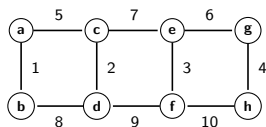
Graph settings

- (G, w) is an edge-weighted graph
 - which can be a pixel-adjacency graph
 - edge weights can represent a gradient of intensity



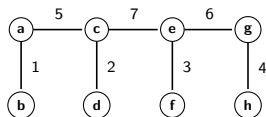
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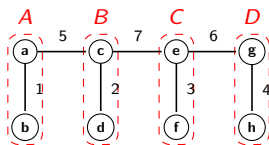
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- For simplification, we consider that *the edges of G have pairwise distinct weights for w , which implies that (G, w) has a single minimum spanning tree (MST)*
- $\mathcal{S} = (M_0, \dots, M_\ell)$ is any sequence of the minima of w ranked by importance according to some regional attribute like extinction values

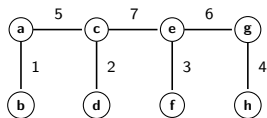


Hierarchical watersheds

- A *hierarchical watershed* of (G, w) for \mathcal{S} is a hierarchy of partitions $(\mathbf{P}_0, \dots, \mathbf{P}_\ell)$ such that, for any $i \in \{0, \dots, \ell\}$:
 - \mathbf{P}_i is the connected component partition of a minimum spanning forest rooted in the minima ranked after i

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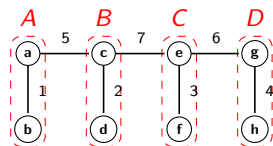
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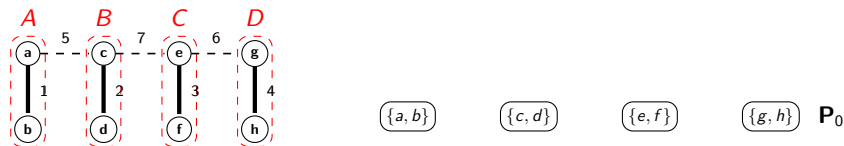


Minima of w

$$\mathcal{S} = (A, C, B, D)$$

Hierarchical watersheds

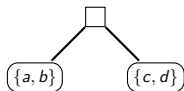
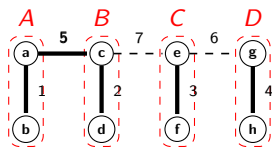
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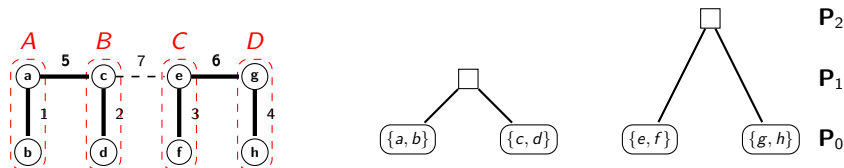
\mathbf{P}_1

\mathbf{P}_0

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Hierarchical watersheds

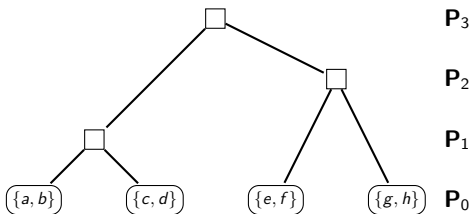
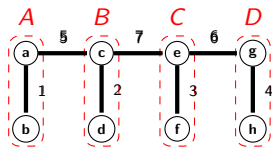
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$(\mathbf{P}_0, \mathbf{P}_1, \mathbf{P}_2, \mathbf{P}_3)$:
hierarchical watershed of (G, w) by $\mathcal{M} = (A, C, B, D)$

Hierarchical watersheds

Definition

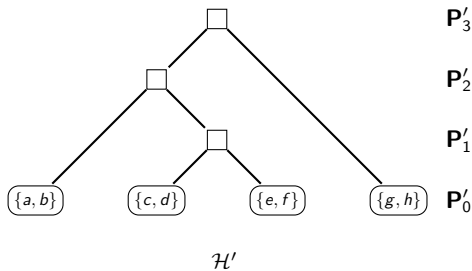
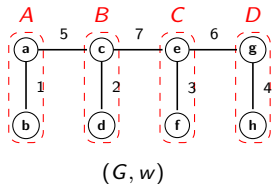
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Remark 1: there are hierarchies which are not hierarchical watersheds of (G, w) , e.g., $(\mathbf{P}'_0, \mathbf{P}'_1, \mathbf{P}'_2, \mathbf{P}'_3)$ is **not** a hierarchical watershed of (G, w)



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Remark 2: a hierarchy \mathcal{H} can be the hierarchical watershed of (G, w) for several sequences of minima of w . An algorithm to count the number of sequences of minima associated to a hierarchical watershed is provided in our companion paper “On the probabilities of hierarchical watersheds” (poster section)

Outlines

- 1 Hierarchical watersheds
- 2 Watersheding transform
- 3 Conclusion and perspectives

Watershed: general idea

- Intuitive idea: invert the algorithm to compute hierarchical watersheds introduced in [C13, N13]

Watershedding algorithm

Input: *weighted graph (G, w) and saliency map s of a hierarchy*

Output: *a saliency map $\omega(s)$, which is the watershedding of s*

- 1 *compute the binary partition hierarchy (by altitude ordering) of w*
- 2 *compute the approximated extinction value (map) ξ_s for s*
- 3 *compute the estimated sequence of minima S_s for s*
- 4 *return the saliency map of the hierarchical watershed of (G, w) for S_s*

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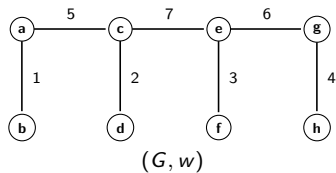
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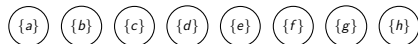
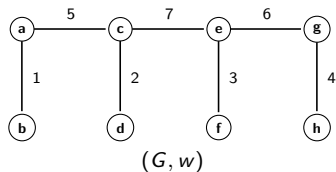
Binary partition hierarchy

- ⇒ Watershedding transform:
1. binary partition hierarchy
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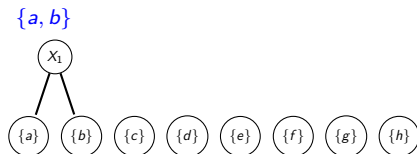
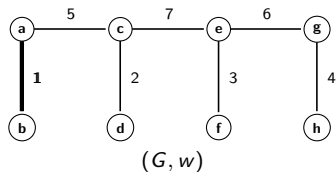
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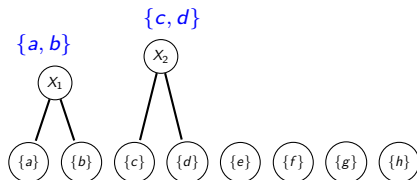
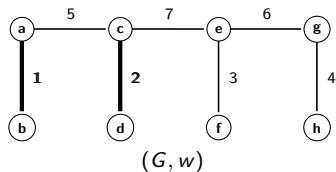
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Any non-leaf region R of \mathcal{B} can be mapped to an edge u_R of G which is then called the building edge of R (shown in blue)

Binary partition hierarchy

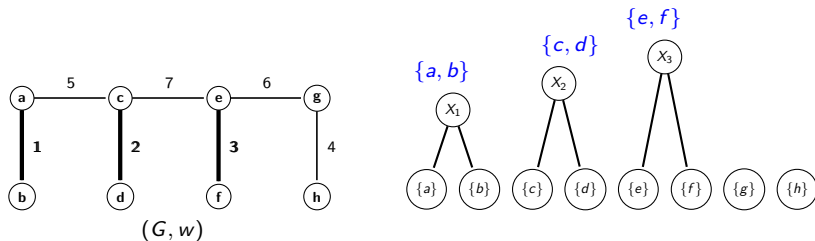
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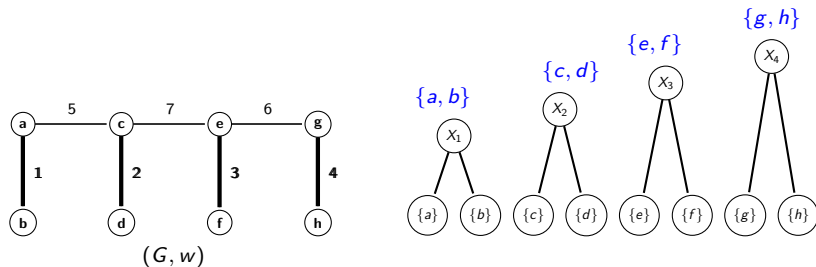
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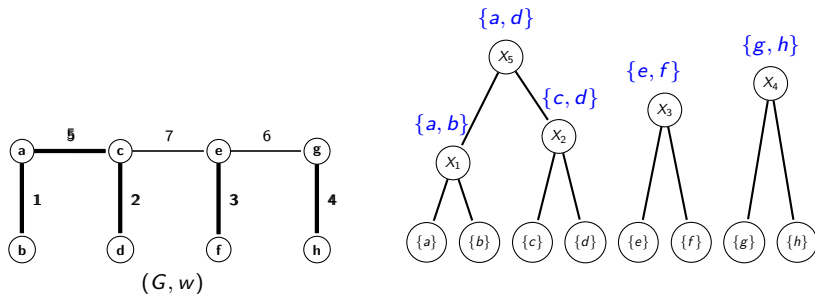
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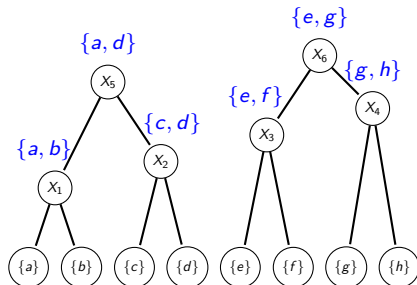
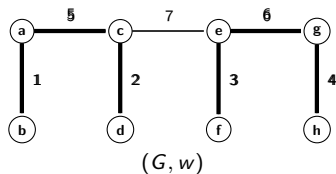
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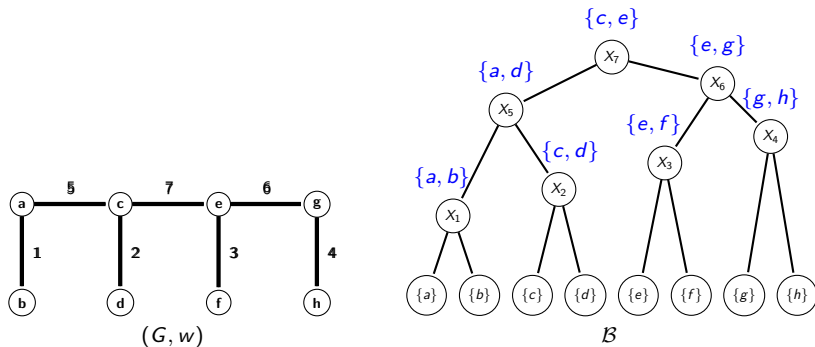
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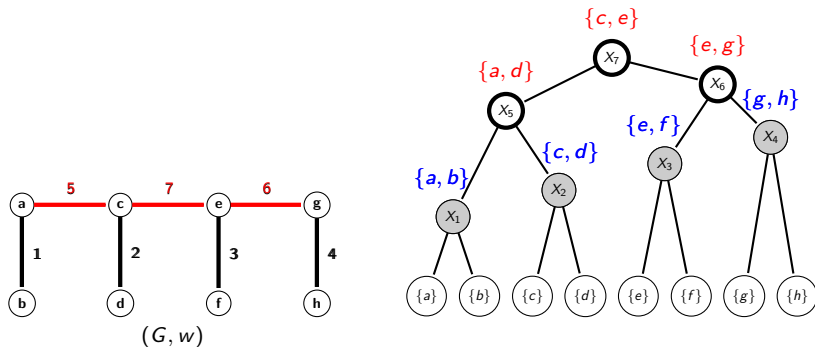
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Any non-leaf region R of \mathcal{B} can be mapped to an edge u_R of G which is then called the building edge of R (shown in blue)
Minima (in grey) and watershed-cut edges (in red)

Approximated extinction map (intuitive idea)

Watershedding transform:
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Definition (Extinction map for a sequence of minima)

Given a sequence $\mathcal{S} = (M_1, \dots, M_n)$ of minima of w , **the extinction map for \mathcal{S}** is a map ϵ such that, for each region R of \mathcal{B} :

- $\epsilon(R)$ is zero if there is no minimum of w included in R
- otherwise, $\epsilon(R)$ is the maximum i such that $M_i \subseteq R$

As established in [C13], the saliency map of a hierarchical watershed for \mathcal{S} can be efficiently computed from the extinction map for \mathcal{S} by one pass on \mathcal{B} .

Definition (Estimated extinction map problem)

Given the saliency map s of a hierarchy \mathcal{H} , find a map ϵ_s , called an approximated extinction map of s such that, if \mathcal{H} is a hierarchical watershed of (G, w) , then this map is an extinction map which induces the hierarchical watershed \mathcal{H} .

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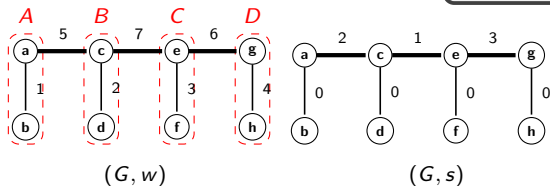
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Approximated extinction map

■ Dominant regions:

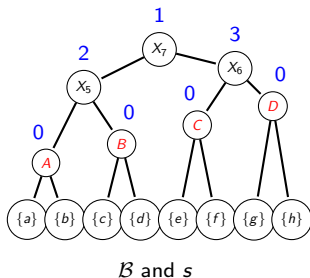
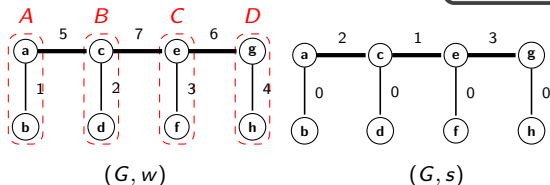
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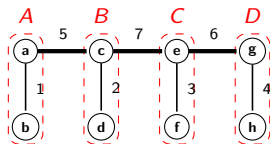
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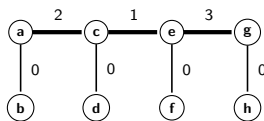
Approximated extinction map

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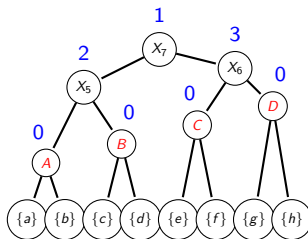
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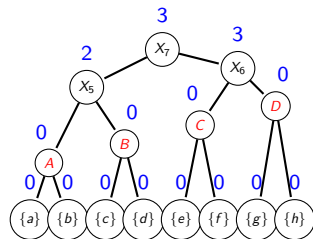
(G, w)



(G, s)



\mathcal{B} and s



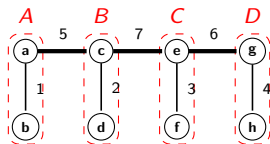
\vee^s

\vee^s : supremum weights among the regions included in each region of \mathcal{B}

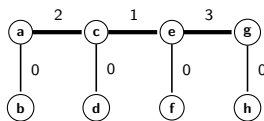
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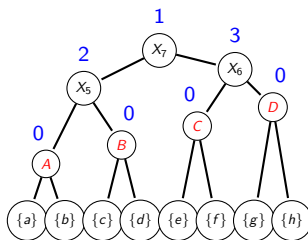
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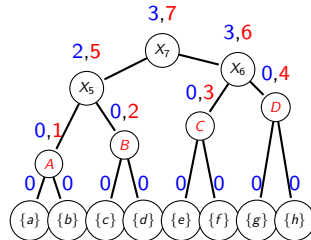
(G, w)



(G, s)



\mathcal{B} and s



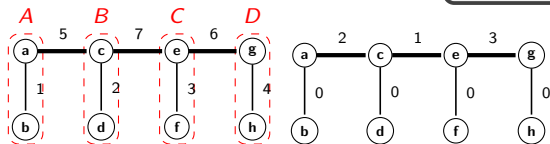
\mathcal{V}^s and w

\mathcal{V}^s : supremum weights among the regions included in each region of \mathcal{B}

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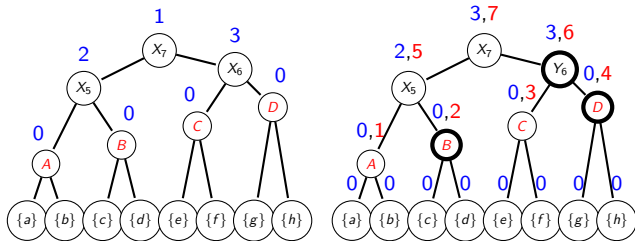
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 - ⇒ 2. approximated extinction map
 3. estimated sequence of minima



(G, w)

(G, s)



\mathcal{B} and s

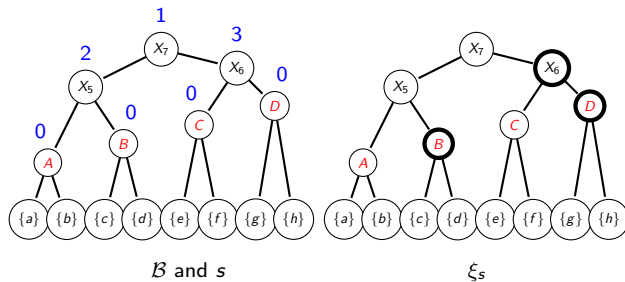
V^s and w

V^s : suppreum weights among the regions included in each region of \mathcal{B}

Dominant regions of \mathcal{B} for s : B , D and Y_6

Approximated extinction map

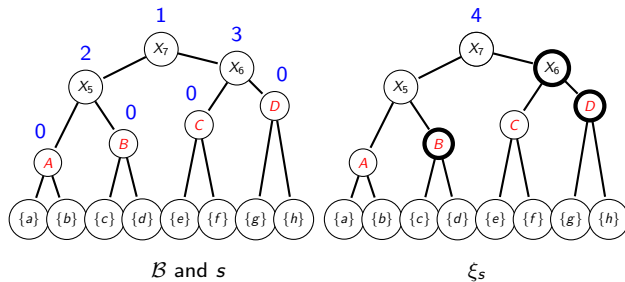
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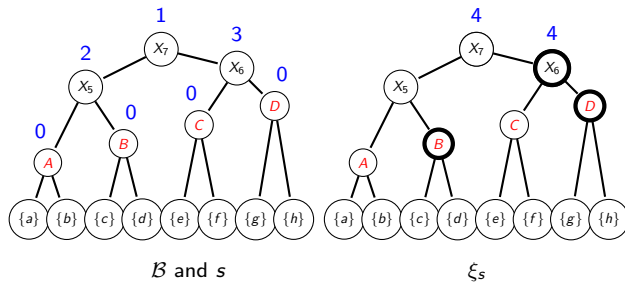
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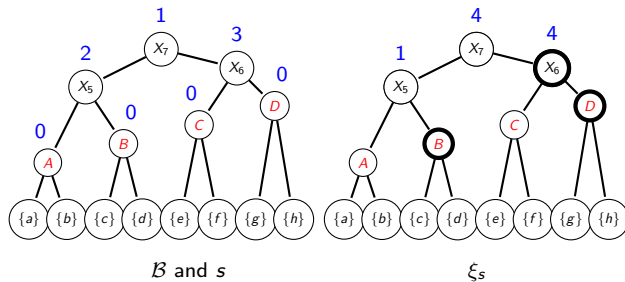
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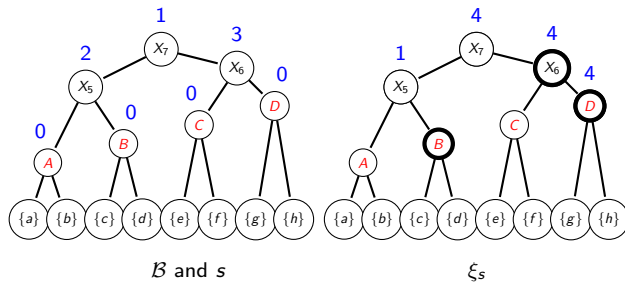
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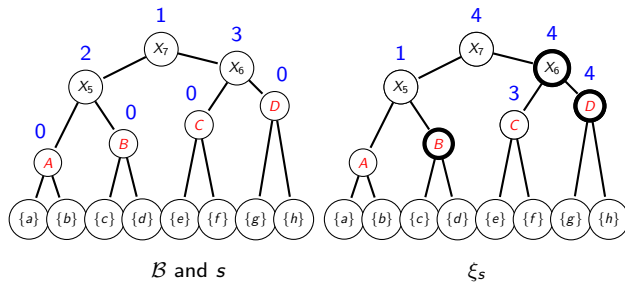
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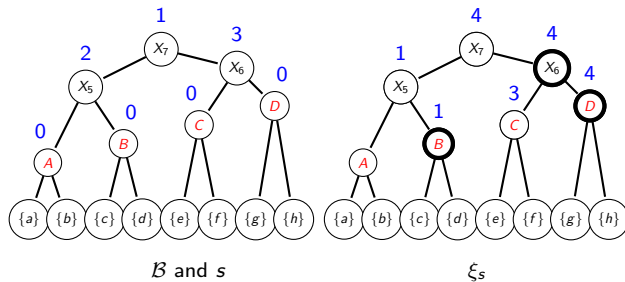
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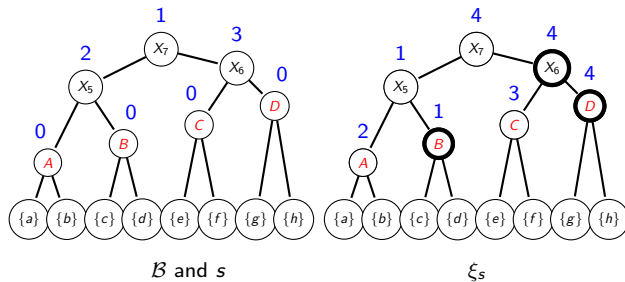
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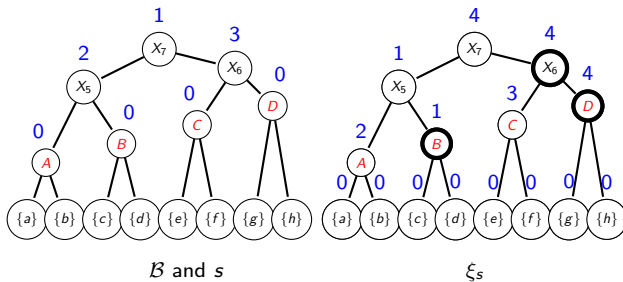
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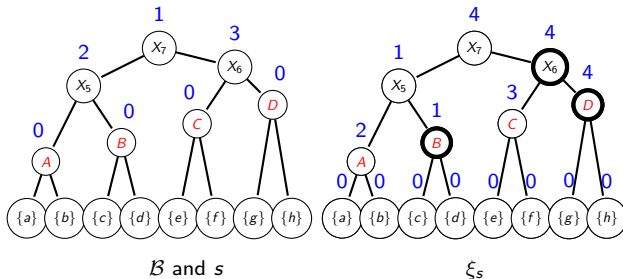
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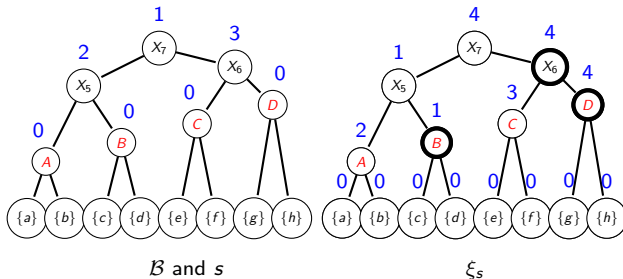
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Estimated sequence of minima for s : $S_f = (B, A, C, D)$

Definition (watershedding)

Let f be a map from the set of edges E into \mathbb{R} and let \mathcal{S} be the estimated sequence of minima of f and let ϵ be the extinction map for \mathcal{S} . The *watershedding of f* is the map $\omega(f)$ from E into \mathbb{R} such that, for any edge u :

- $\omega(f)(u) = \min\{\epsilon(R) \mid R \text{ is a child of } R_u\},$

where R_u is the region of \mathcal{B} whose building edge is u .

Watershedding operator: properties

Property (Idempotence)

Let f be a map from E into \mathbb{R} . The watershedding $\omega(\omega(f))$ of $\omega(f)$ is equal to the watershedding $\omega(f)$ of f .

Property (Invariance domain)

Let \mathcal{H} be a hierarchy and let f be the saliency map of \mathcal{H} . The watershedding of f is equal to f if and only if \mathcal{H} is a hierarchical watershed of (G, w) .

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Outlines

- 1 Hierarchical watersheds
- 2 Watersheding transform
- 3 Conclusion and perspectives**

Summary

- *Introduction of an idempotent operator which converts any hierarchy into a hierarchical watersheds*

Perspectives

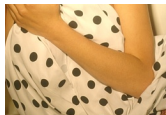
- *Extension to arbitrary weighted graphs*
- *Open question: does the watershed optimization any objective function?*

Watershedding operator: applications

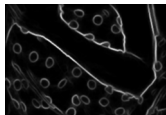
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 - e.g., circularity

Watershedding operator: applications

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I



G



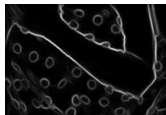
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Watershedding operator: applications

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I



G



\mathcal{H}_c



\mathcal{H}_{cc}

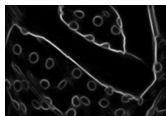
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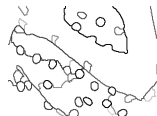
G



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\mathcal{H}_{cc}

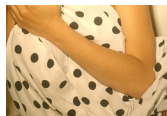


\mathcal{H}_w

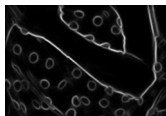
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\mathcal{H}_{cc}



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I'



G'

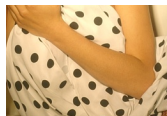


\mathcal{H}'_c

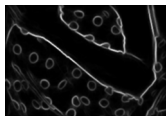
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I'



G'



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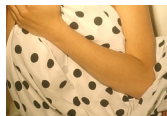


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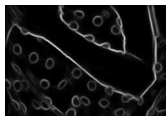
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